Interval minors of binary matrices

Stanislav Kučera

Computer Science Institute, Charles University

Joint work with Vít Jelínek.

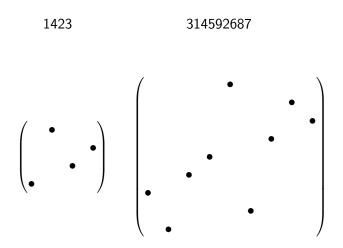
June 26, 2017

Stanislav Kučera (Charles Univ.)

Interval minors

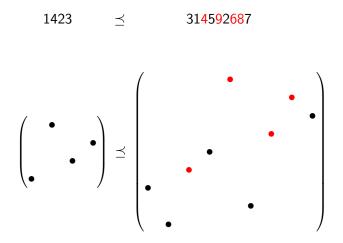
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From permutations to binary matrices



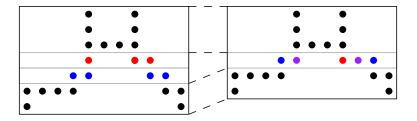
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From permutations to binary matrices



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Row contraction



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Definition 1

A binary matrix P is an interval minor of a binary matrix M, if P can be obtained from M by a sequence of:

- line (row or column) contractions, and
- one-entry deletions.

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Definition 2

A binary matrix $P \in \{0,1\}^{k \times l}$ is an interval minor of a binary matrix M, if there is a partitioning of M by k-1 horizontal and l-1 vertical lines to $k \cdot l$ rectangular submatrices M[i,j], $1 \le i \le k$ and $1 \le j \le l$ such that if $P_{i,j}$ is a one-entry then M[i,j] contains a one-entry.

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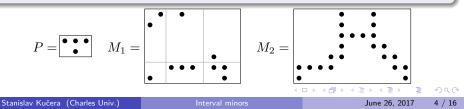
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Definition

A binary matrix M is P-critical for a binary matrix P, if P is not an interval minor of M and a change of an arbitrary zero-entry to a one-entry creates a matrix containing an interval minor P.

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- *P* with two one-entries:

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Rotation does not break the property.

• *P* with three one-entries:

• *P* with four one-entries:

$$P = (\bullet, \bullet) \qquad M = \left(\bullet, \bullet, \bullet \bullet \right) \text{ or } \left(\bullet, \bullet, \bullet \bullet \right)$$
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What do they all have in common?

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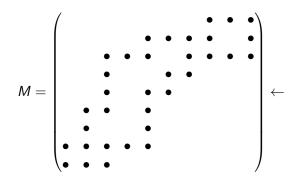
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Question 1

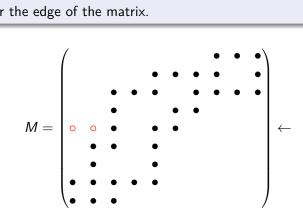
Is this true for every matrix P?

Definition

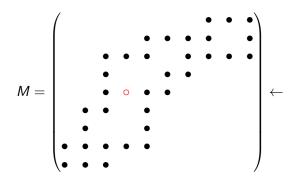
Definition



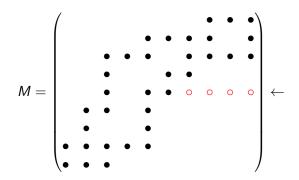
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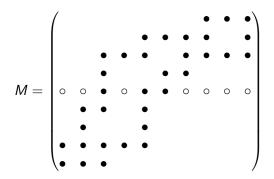
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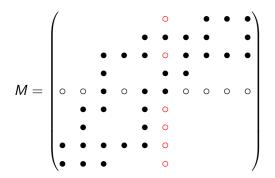
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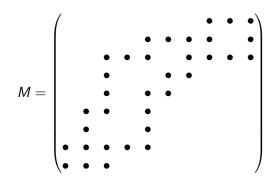
A zero-interval in a matrix M is an interval (contiguous sequence) of zero-entries in a row or column of M, bounded from each side by a one-entry or the edge of the matrix.

Definition

The row complexity of a matrix M is the maximum number of zero-intervals in a row of M. The column complexity of a matrix M is the maximum number of zero-intervals in a column of M. The complexity of a matrix M is the maximum of its row and column complexity.

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Definition



Bounding matrices

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Is every matrix *P* bounding?

No!

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Bounding matrices – theorem

Theorem 1

A binary matrix P is non-bounding if and only if any rotation of $P_1 = \begin{pmatrix} \bullet \\ \bullet \end{pmatrix}$ is an interval minor of P.

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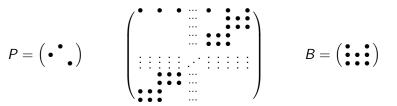
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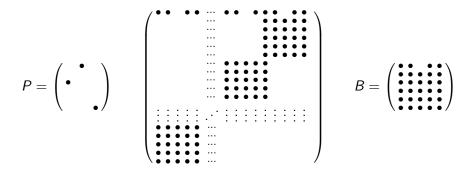


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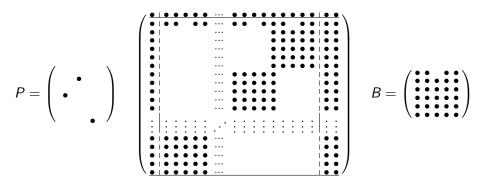


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If a matrix P is bounding then every P' interval minor of P is bounding.

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Definition

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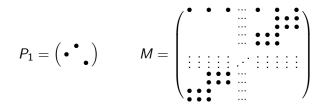
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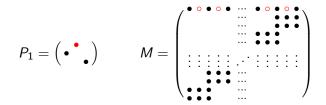
Corollary 2

A matrix P is row-bounding if and only if P is column-bounding.

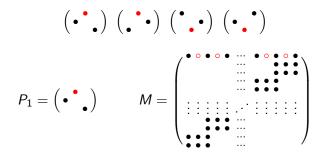
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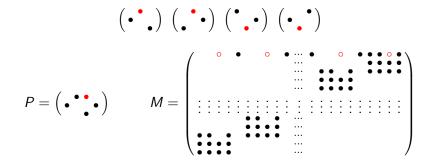


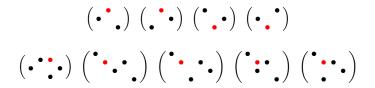
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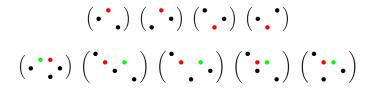
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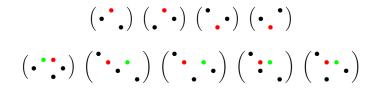




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Question

Does every difficult one-entry share the row with a trivially difficult one-entry?

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- Can Corollary 1 or Corollary 2 be proved without using Theorem 1?
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For a bounding matrix P let f(P) be the minimum $k \in \mathbb{N}$ such that the complexity of every P-critical matrix is less than or equal to k.

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For a bounding matrix P let f(P) be the minimum $k \in \mathbb{N}$ such that the complexity of every P-critical matrix is less than or equal to k.

• What is the magnitude of f(P) for a bounding binary matrix P?

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- What is the magnitude of f(P) for a bounding binary matrix P? Is it linear? Quadratic? ...
- Is f(P) linear for P of size $2 \times n$?
- Let P be an interval minor of a bounding binary matrix P'. Is f(P) ≤ f(P')?

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Thank you.

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