Automedians sets of permutation: extended abstract *

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Abstract. Given a set $\mathcal{A} \subseteq \mathbb{S}_n$ of m permutations of $\{1, 2, \ldots, n\}$ and a distance function d, the **median** problem consists of finding the set $\mathcal{M}(\mathcal{A})$ of all the permutations that are the "closest" of the m given permutations. Here, we study the automedian case of the problem, i.e. when $\mathcal{A} = \mathcal{M}(\mathcal{A})$, under the Kendall- τ distance which counts the number of order disagreements between pairs of elements of permutations.

1 Introduction

The problem of finding medians of a set of m permutations of $\{1, 2, \ldots, n\}$ under the Kendall- τ distance [8, 12] is often cited in the literature as the Kemeny Score Problem [7]. Giving an ordering of n political candidate by m voters, the problem consists of finding a *Kemeny consensus*: an order of the candidates that agrees the most with the order of the m voters. This problem is polynomial-time solvable for m = 2, has been proved to be NP-complete when $m \geq 4$, m even (first proved in [4], then corrected in [3]), but its complexity remains unknown for $m \geq 3$, m odd. Some approximate and fixed parameters algorithms have been derived [1, 2, 6, 9-11, 13]. Here we investigate automedian sets of permutations (when the sets of permutations \mathcal{A} is equal to the set of its medians $M(\mathcal{A})$) by building on a previously published extended abstract [5] and constructing new variants of automedian sets of permutations using some combinatorial properties of those sets. Note that all proofs are omitted here but will be included in the full version of this article.

2 Median of permutation

2.1 Definitions and notations

A **permutation** π is a bijection of $[n] = \{1, 2, ..., n\}$ onto itself. The set of all permutations of [n] is denoted \mathbb{S}_n . As usual we denote a permutation π of [n] as $\pi = \pi_1 \pi_2 ... \pi_n$. The **identity permutation** corresponds to the identity bijection of [n] and is denoted i = 12...n. Given two permutations π and $\sigma \in \mathbb{S}_n$, $\pi\sigma$ (or $\pi \circ \sigma$) denote the usual **product of permutation** *i.e.* their composition as functions. Let $\mathcal{A} \subseteq \mathbb{S}_n$ be a set of permutations of [n]; we will denote its **cardinality** by $\#\mathcal{A}$.

The **Kendall-** τ **distance**, denoted d_{KT} , counts the number of order disagreements between pairs of elements of two permutations and can be defined formally as follows: for permutations π and σ of [n], we have that $d_{KT}(\pi, \sigma) = \#\{(i, j) | i < j \text{ and } [(\pi_i^{-1} < \pi_j^{-1} \text{ and } \sigma_i^{-1} > \sigma_j^{-1}) \text{ or } (\pi_i - 1 > \pi_j - 1 \text{ and } \sigma_i^{-1} < \sigma_j^{-1})]\}$. Given any set of permutations $\mathcal{A} \subseteq \mathbb{S}_n$ and a permutation π , we have $d_{KT}(\pi, \mathcal{A}) = \sum_{\sigma \in \mathcal{A}} d_{KT}(\pi, \sigma)$.

The problem of finding a median of a set of permutations \mathcal{A} under the **Kendall**- τ distance is the following. Given $\mathcal{A} \subseteq \mathbb{S}_n$, we want to find a permutation π^*

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of [n] such that $d_{KT}(\pi^*, \mathcal{A}) \leq d_{KT}(\pi, \mathcal{A}), \forall \pi \in \mathbb{S}_n$. Finally, let us define $\mathcal{M}(\mathcal{A})$ as **the** set of all medians of \mathcal{A} i.e. $\mathcal{M}(\mathcal{A}) = \{\pi \in \mathbb{S}_n | \forall \sigma \in \mathbb{S}_n, d_{KT}(\pi, \mathcal{A}) \leq d_{KT}(\sigma, \mathcal{A})\}$. In this article, we are interested to investigate automedian sets of permutations i.e sets of permutations, \mathcal{A} , that are equal to their set of medians, $\mathcal{M}(\mathcal{A})$.

2.2 Previous work on automedian sets

In a previous extended abstract [5], we gave two special simple cases of automedian sets. The first case (see Theorem 1) is when \mathcal{A} is composed of a permutation and its "circular rotations". The second one (see Theorem 2) is when \mathcal{A} consist of permutations composed of a common S_k kernel with some (or none) shared fixed points.

Theorem 1 ([5]). Let $\pi = \pi_1 \dots \pi_n$ be a permutation of [n] and let $\uparrow_i \pi = \pi_{i+1} \dots \pi_n \pi_1 \dots \pi_i$. If $\mathcal{A} = \{\uparrow_i \pi \mid 0 \leq i \leq n-1\}$ then $\mathcal{A} = \mathcal{M}(\mathcal{A})$ and, $\forall i, 0 \leq i \leq n-1, d_{KT}(\uparrow_i \pi, \mathcal{A}) = \frac{(n+1)n(n-1)}{6}$.

Theorem 2 ([5]). If $\mathcal{A} = \{\pi = 1 2 \dots \ell \sigma \ \ell + k + 1 \dots n \mid \sigma \text{ is any permutation of } \{\ell + 1, \dots, \ell + k\}\}$, for any $0 \leq \ell \leq n - 1$, $1 \leq k \leq n$, then $\mathcal{A} = \mathcal{M}(\mathcal{A})$ and for $\pi \in \mathcal{A}$, $d_{KT}(\pi, \mathcal{A}) = \frac{k(k-1)k!}{4}$. Note that $\mathcal{A} = \mathbb{S}_n$, is the special case when $\ell = 0$ and k = n.

3 Group action and morphism

Let $\mathcal{P}(\mathbb{S}_n)$ be the set of all subsets of \mathbb{S}_n . Composition of permutations turns \mathbb{S}_n into a group and we get a **left group action** on $\mathcal{P}(\mathbb{S}_n)$:

$$\begin{array}{ccc} \circ : \mathbb{S}_n \times \mathcal{P}(\mathbb{S}_n) \longrightarrow \mathcal{P}(\mathbb{S}_n) \\ (\pi \ , \ \mathcal{A}) & \mapsto & \pi \circ \mathcal{A}, \end{array}$$

by setting $\pi \circ \mathcal{A} = \{\pi \circ \sigma \mid \sigma \in \mathcal{A}\}$. We will first show two small lemmas about the invariability of the Kendall- τ distance under this action. Note that we denote $\pi \circ \sigma$ by $\pi \sigma$ in the rest of this article.

Lemma 1. Let π, σ, ψ be permutations of \mathbb{S}_n . Then $d_{KT}(\pi\sigma, \pi\psi) = d_{KT}(\sigma, \psi)$.

Lemma 2. Let $\mathcal{A} \subseteq \mathbb{S}_n$ be a set of permutations and let π be any permutation of \mathbb{S}_n . Then, $d_{KT}(\pi\sigma, \pi A) = d_{KT}(\sigma, A)$.

We now have everything in place to show the following:

Theorem 3. Let $\mathcal{M} : \mathcal{P}(\mathbb{S}_n) \longrightarrow \mathcal{P}(\mathbb{S}_n)$ be the function that maps a set of permutations $\mathcal{A} \subseteq \mathbb{S}_n$ to the set of its medians $\mathcal{M}(\mathcal{A})$. The function \mathcal{M} is a morphism of action i.e. $\pi \mathcal{M}(\mathcal{A}) = \mathcal{M}(\pi \mathcal{A})$.

Recall that the **orbit** of \mathcal{A} under the action \circ is the set of all $\sigma \mathcal{A}, \sigma \in \mathbb{S}_n$. We have the following first result on automedian sets:

Corollary 1. Let \mathcal{A} and \mathcal{B} be two sets of permutations of \mathbb{S}_n in the same orbit of the action $\circ : \mathbb{S}_n \times \mathcal{P}(\mathbb{S}_n) \to \mathcal{P}(\mathbb{S}_n)$. Then $\mathcal{B} = \mathcal{M}(\mathcal{B}) \iff \mathcal{A} = \mathcal{M}(\mathcal{A})$.

4 Decomposable permutations and direct sum

Now that we have derived some basic properties of the sets of medians, let us explore ways to built new automedian sets of decomposable permutations, using the direct sum of already known sets. First, let us recall some definitions and set some notations.

4.1 Definitions and notations

Definition 1. A permutation $\pi \in \mathbb{S}_n$ will be called k-decomposable, $1 \le k \le n-1$, if $i > k \Leftrightarrow \pi_i > k, \forall i \in \{1, 2, ..., n\}$.

Definition 2. A permutation $\pi \in S_n$ will be called indecomposable, if it is not k-decomposable for any $k \in \{1, 2, ..., n-1\}$.

Definition 3. A set of permutations $\mathcal{A} \subseteq \mathbb{S}_n$ will be called k-decomposable, $1 \leq k \leq n-1$, if all of its permutations are k-decomposable.

Definition 4. Let $\mathcal{A} \subseteq \mathbb{S}_n$ be a k-decomposable set, $1 \leq k \leq n-1$, and let $\pi \in \mathbb{S}_n$ be any permutation. The set $\pi \mathcal{A}$ is called **k-separable**. If \mathcal{A} is not **k-separable** for any k, then it we be called **inseparable**.

Note that to lighten notation, if there are many strictly ascending values $k_1, k_2..., k_m$ such that a permutation or a set of permutations is k_i -decomposable (resp. k_i -separable), $1 \le i \le m$, we will say that it is $k_1; k_2; ..., k_m$ -decomposable (resp. $k_1; k_2; ..., k_m$ -separable).

Example 1. Let $\pi = 32145$, then π is 3;4-decomposable. Let $\sigma = 54321$, then σ is indecomposable.

Definition 5. Let $\pi \in S_k$ and $\sigma \in S_\ell$ be two permutations of length k and ℓ , respectively. The **direct sum** of π and σ , denoted $\pi \oplus \sigma$, is the permutation of length $k + \ell$ defined as $\pi \oplus \sigma = \pi_1 \pi_2 \dots \pi_k (\sigma_1 + k) (\sigma_2 + k) \dots (\sigma_\ell + k)$.

Example 2. Let $\pi = 132$ and $\sigma = 2143$, then $\pi \oplus \sigma = 1325476$.

Note: A permutation $\pi \in S_n$ that is $k_1, k_2..., k_m$ -decomposable can be written as the direct sum of m + 1 indecomposable permutations of length $k_i - k_{i-1}$, for $1 \le i \le m + 1$, with $k_0 = 0$ and $k_{m+1} = n$.

Example 3. The permutation $\pi = 32145 \in \mathbb{S}_5$ is 3;4-decomposable. It can be written as $\pi = 321 \oplus 1 \oplus 1$, where $321 \in \mathbb{S}_3$ and $1 \in \mathbb{S}_1$ are indecomposable permutations.

We can extend the definition of direct sum to sets of permutations in the following way:

Definition 6. Let $\mathcal{A} \subseteq \mathbb{S}_k$ and $\mathcal{B} \subseteq \mathbb{S}_\ell$ be two sets of permutations of length k and ℓ , respectively. The **direct sum** of \mathcal{A} and \mathcal{B} , denoted $\mathcal{A} \oplus \mathcal{B}$, is defined as $\mathcal{A} \oplus \mathcal{B} = \{\pi \oplus \sigma \mid \pi \in \mathcal{A} \text{ and } \sigma \in \mathcal{B}\}.$

Note that by definition, $\mathcal{A} \oplus \mathcal{B}$ is k-decomposable.

Example 4. Let $\mathcal{A} = \{132, 312\}$ and $\mathcal{B} = \{2143, 2314, 2431\}$. Then $\mathcal{A} \oplus \mathcal{B}$ is 3-decomposable and $\mathcal{A} \oplus \mathcal{B} = \{1325476, 1325647, 1325764, 3125476, 3125647, 3125764\}$. Let $\pi = 3271653$, then $\pi(\mathcal{A} \oplus \mathcal{B}) = \{3726145, 3726451, 3726514, 7326145, 7326451, 7326514\}$ is 3-separable.

It is easy to show that the direct sum \oplus is an associative operation and so we have, for any permutations π, σ and ϕ , that $(\pi \oplus \sigma) \oplus \phi = \pi \oplus (\sigma \oplus \phi)$. This associativity easily extends to the direct sums of sets of permutations.

4.2 Theorems and properties

Let us first show this easy result about the Kendall- τ distance of permutations that are direct sums of smaller permutations.

Lemma 3. Giving ρ and ψ , two permutations of \mathbb{S}_k , ϕ and θ , two permutations of \mathbb{S}_ℓ , let π and σ be the two permutations of \mathbb{S}_n , $n = k + \ell$, such that $\pi = \rho \oplus \phi$ and $\sigma = \psi \oplus \theta$. Then, $d_{KT}(\pi, \sigma) = d_{KT}(\rho, \psi) + d_{KT}(\phi, \theta)$.

Now, let us explore some properties connecting the direct sums of sets of permutations and the median of those sets.

Theorem 4. Let $\mathcal{A} \subseteq \mathbb{S}_k$ and $\mathcal{B} \subseteq \mathbb{S}_\ell$, $k + \ell = n$, be two sets of permutations. Then, $\mathcal{M}(\mathcal{A} \oplus \mathcal{B}) = \mathcal{M}(\mathcal{A}) \oplus \mathcal{M}(\mathcal{B}).$

We can use Theorem 4 and the next Lemma to obtain new automedian sets by taking the direct sums of already known ones. This implies that automedian sets of permutations are closed under the operation direct sum.

Lemma 4. Let $\mathcal{A}, \mathcal{C} \subseteq \mathbb{S}_k$ and $\mathcal{B}, \mathcal{D} \subseteq \mathbb{S}_\ell$ be four sets of permutations. If $\mathcal{A} \oplus \mathcal{B} = \mathcal{C} \oplus \mathcal{D}$, then $\mathcal{A} = \mathcal{C}$ and $\mathcal{B} = \mathcal{D}$.

Corollary 2. Let $\mathcal{A} \subseteq \mathbb{S}_k$ and $\mathcal{B} \subseteq \mathbb{S}_\ell$ be two sets of permutations. Then $\mathcal{A} = \mathcal{M}(\mathcal{A})$ and $\mathcal{B} = \mathcal{M}(\mathcal{B}) \iff \mathcal{A} \oplus \mathcal{B} = \mathcal{M}(\mathcal{A} \oplus \mathcal{B})$

Now, we have everything in place to show that any k-decomposable automedian set of permutations comes from the direct sums of automedian sets of smaller permutations.

Theorem 5. Let $C \subset \mathbb{S}_n$ be a k-decomposable automedian set. Then there exist automedian sets of permutations $\mathcal{A} \subseteq \mathbb{S}_k$ and $\mathcal{B} \subseteq \mathbb{S}_\ell$, $k + \ell = n$, such that $\mathcal{A} \oplus \mathcal{B} = C$.

Corollary 3. Let $C \subset \mathbb{S}_n$ be a k-separable automedian set. Then there exist $\pi \in \mathbb{S}_n$ and automedian sets of permutations $A \subseteq \mathbb{S}_k$ and $B \subseteq \mathbb{S}_\ell$, $k + \ell = n$, such that $\pi(A \oplus B) = C$.

Lemma 5. Let $C \subseteq \mathbb{S}_n$ be a $k_1; k_2; \ldots; k_m$ -decomposable automedian set. Then there exist automedian sets of permutations $\mathcal{A}_1 \subseteq \mathbb{S}_{k_1}, \mathcal{A}_2 \subseteq \mathbb{S}_{k_2-k_1}, \ldots, \mathcal{A}_m \subseteq \mathbb{S}_{k_m-k_{m-1}}, \mathcal{A}_{m+1} \subseteq \mathbb{S}_{n-k_m}$, such that $\mathcal{A}_1 \oplus \mathcal{A}_2 \oplus \ldots \oplus \mathcal{A}_m \oplus \mathcal{A}_{m+1} = C$.

Corollary 4. Let $C \subset \mathbb{S}_n$ be a $k_1; k_2; \ldots; k_m$ -separable automedian set. Then there exist $\pi \in \mathbb{S}_n$ and automedian sets of permutations $\mathcal{A}_1 \subseteq \mathbb{S}_{k_1}, \mathcal{A}_2 \subseteq \mathbb{S}_{k_2-k_1}, \ldots, \mathcal{A}_m \subseteq \mathbb{S}_{k_m-k_{m-1}}, \mathcal{A}_{m+1} \subseteq \mathbb{S}_{n-k_m}$, such that $\pi(\mathcal{A}_1 \oplus \mathcal{A}_2 \oplus \ldots \oplus \mathcal{A}_m \oplus \mathcal{A}_{m+1}) = C$.

4.3 Counting automedian sets

Corollary 4 shows us that any separable automedian set of permutations $\subseteq S_n$ can be written as a composition of a permutation $\pi \in S_n$ with the direct sum of inseparable automedian sets. It follows that to solve the problem of finding all automedian sets of permutations of S_n , we only need to find inseparable ones because we can get all the others by combining smaller inseparable sets. To count all the automedian sets of S_n , let \mathcal{AM}_n be the set of all automedian sets of S_n , i.e. $\mathcal{AM}_n = \{\mathcal{A} = \mathcal{M}(\mathcal{A}) \mid \mathcal{A} \subseteq S_n\}$ and let \mathcal{I}_n be the set of all inseparable ones, i.e., $\mathcal{I}_n = \{\mathcal{A} = \mathcal{M}(\mathcal{A}) \mid \mathcal{A} \subseteq S_n \text{ and } \mathcal{A} \text{ inseparable}\}.$ Now, an automedian set of permutations $\mathcal{A} \subseteq \mathbb{S}_n$ is either inseparable, and there are $\#\mathcal{I}_n$ of those sets or it is separable. In that case, let $i, 1 \leq i \leq n-1$ be the smallest value for which \mathcal{A} is *i*-separable. Everything we have built so far gives us that $\#\mathcal{A}\mathcal{M}_n = \#\mathcal{I}_n + \sum_{i=1}^{n-1} {n \choose i} \times \#\mathcal{I}_i \times \#\mathcal{A}\mathcal{M}_{n-i}$, where the ${n \choose i}$ in the formula counts the number of ways to pick the *i* elements among *n* that are in the inseparable part of \mathcal{A} . We still do not know how to generate all inseparable automedian sets but we did count them for small *n* and got the following values of $\#\mathcal{A}\mathcal{M}_n$ and $\#\mathcal{I}_n$:

n	$\#\mathcal{I}_n$	$\# \mathcal{AM}_n$
1	1	1
2	1	3
3	3	15
4	27	117
5	≥ 429	≥ 1389
6	≥ 8889	≥ 23667

5 Conclusion and future works

The median problem has been essentially studied for the permutation group S_n and under the Kendall- τ distance. From a combinatorial point of view, it would be very interesting to generalize the median problem to the signed permutations group \mathcal{B}_n , as well as to other Coxeter groups which are known to behave in similar manner to these two groups for a large number of combinatorial questions.

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